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NUCLEAR TECHNOLOGY
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Subcriticality Measurement using Feynman- α with a Fully Random Sampling and Second- Order Filtering Technique for AGN-201K

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- ❖ Most nuclear facilities are designed to have conservative subcritical margin to prevent accidentally uncontrolled neutron multiplications.
- ❖ Therefore, an accurate real-time measurement of subcriticality can provide a helpful way to guarantee the safe operation of nuclear facilities.
- ❖ Noise analysis methods have been studied for a long time for this purpose.
- ❖ In this work, subcriticality experiment is performed with the **Feynman- α method** at AGN-201K which is zero-power research and training reactor in our country.
- ❖ To reduce computing time and for improve accuracy near five critical states in estimating the prompt neutron decay constant, **a fully random sampling technique coupled with the second order differential filtering is devised** to effectively process the data obtained with a fine gate time within reduced computing time.

❖ *Noise analysis method*

- Noise analysis methods are based on the same basic premise that the properties of a subcritical system can be determined by measuring the fluctuations in the fission chain processes that depend on the stochastic nature of the birth and death of neutrons.
- So, if the time of the source or detection event are measurable, the distribution of the times between the source (or detection) event and detection event would provide a direct indication of the dynamic properties of the subcritical system.

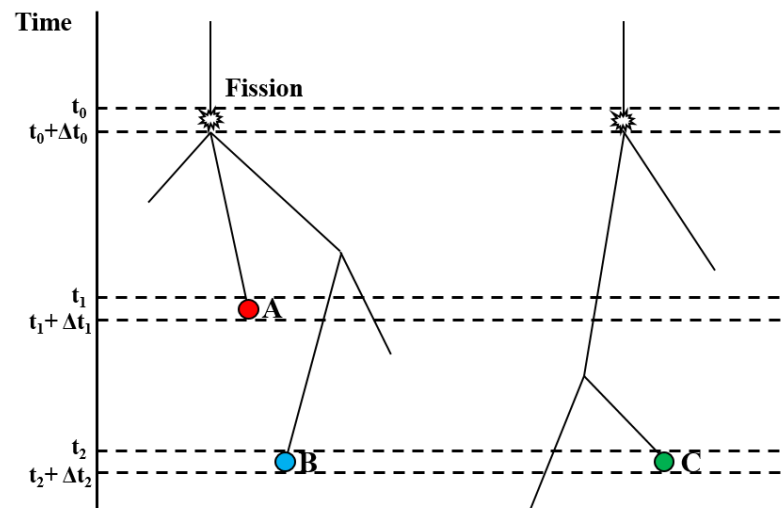


Fig. 1 Random branching process of fission neutron

❖ Feynman- α method

- The Feynman- α method can be derived from the Rossi- α method. This method can determine the prompt neutron decay constant (α) by considering the ratio of the variance to the mean of neutron counts collected in a fixed time interval (i.e., gate time).

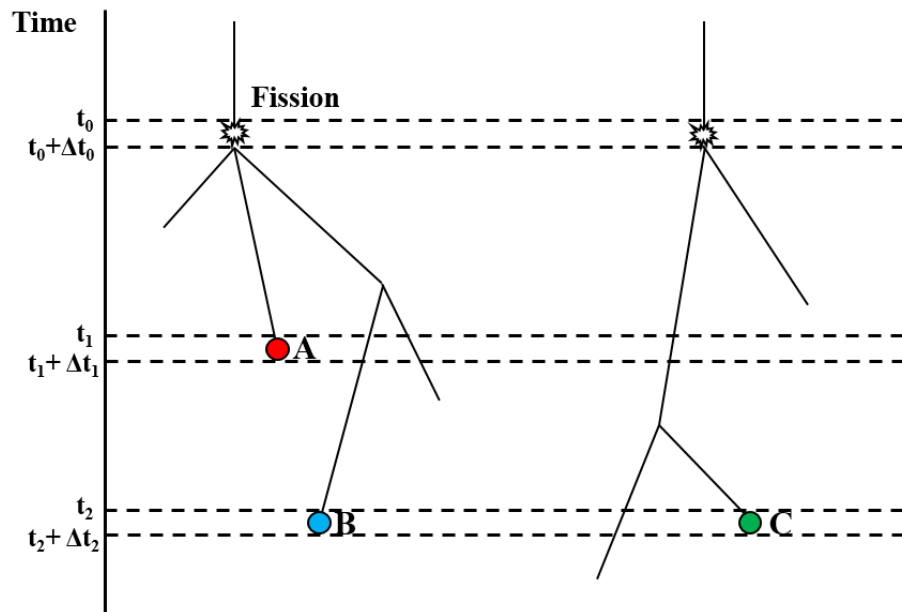


Fig. 1 Random branching process of fission neutron

Rossi- α method

$$p(t_1, t_2) dt_1 dt_2 = F \varepsilon^2 \left(F + \frac{D_v k_p^2}{2(1 - k_p) \Lambda} e^{-\alpha(t_2 - t_1)} \right) dt_1 dt_2$$

Integration of it over time

$$\int_{t_2=0}^{\tau} \int_{t_1=0}^{t_2} p(t_1, t_2) dt_1 dt_2 = 1 + \frac{\varepsilon D_v}{\rho_p^2} \left(1 - \frac{1 - e^{-\alpha\tau}}{\alpha\tau} \right); \quad \rho_p = \frac{k_p - 1}{k_p}$$

Feynman- α method

$$Y(\tau) = \frac{\overline{C_k^2} - (\overline{C_k})^2}{\overline{C_k}} - 1 = Y_{\infty} \left[1 - \frac{1 - e^{-\alpha\tau}}{\alpha\tau} \right]$$

THEORY AND METHOD

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❖ Feynman- α method

$$Y(\tau) = \frac{\overline{C_k^2} - (\overline{C_k})^2}{\overline{C_k}} - 1 = Y_\infty \left[1 - \frac{1 - e^{-\alpha\tau}}{\alpha\tau} \right],$$

- where Y is defined as the variance-to-mean ratio of a series of neutron counts (C_k) with a gate time τ subtracted by 1. The saturated correlation amplitude Y_∞ includes detector efficiency ε , Diven's factor D_v and prompt reactivity ρ_p .

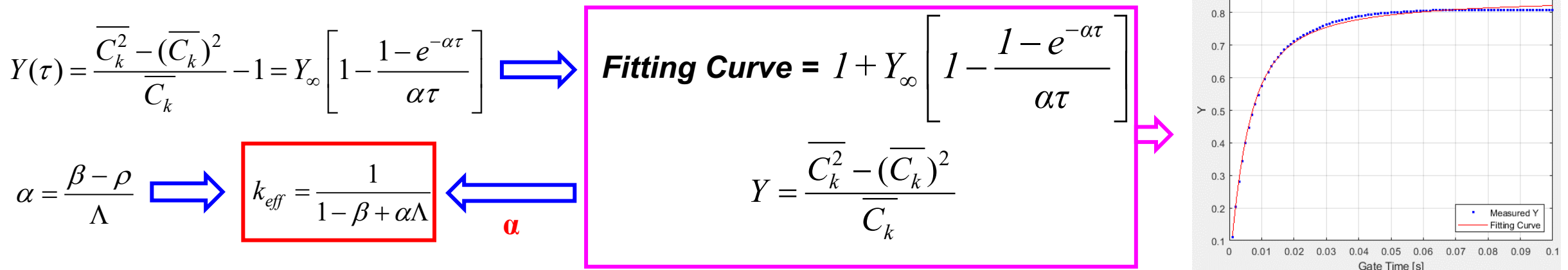


Fig. 2 Determination of prompt neutron decay constant by using Feynman- α fitting curve



❖ 2nd order Feynman- α differential filtering method

- However, the conventional Feynman- α method suffers from **the divergence of the variance near the critical state**.
- To circumvent the divergence of the variance, **Bennett (1960)** proposed an improved method with **differences of the counts between adjacent gates**.
- **Hashimoto et al. (1997)** generalized the Bennett's method to develop a difference-filtering technique and proposed a usage of the higher-order filtering for Feynman- α method to reduce the effect of reactor-power drift during a measurement.

$$\sigma_2(\tau) = \frac{\left(C_k - \frac{1}{2}C_{k-1} - \frac{1}{2}C_{k+1} \right)^2}{\overline{C_k} - \frac{3}{2}}$$
$$= Y_\infty \left(1 - \frac{\frac{5}{3} - \frac{5}{2}e^{-\alpha\tau} + e^{-2\alpha\tau} - \frac{1}{6}e^{-3\alpha\tau}}{\alpha\tau} \right)$$

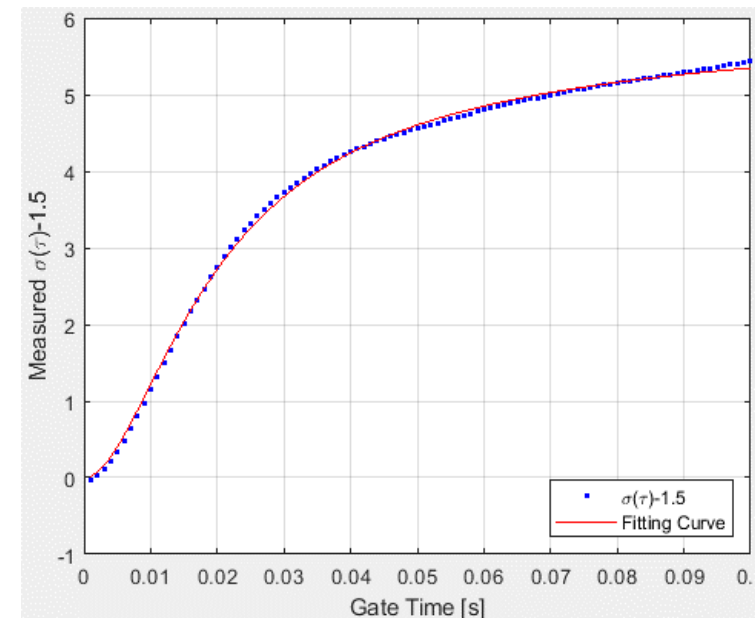
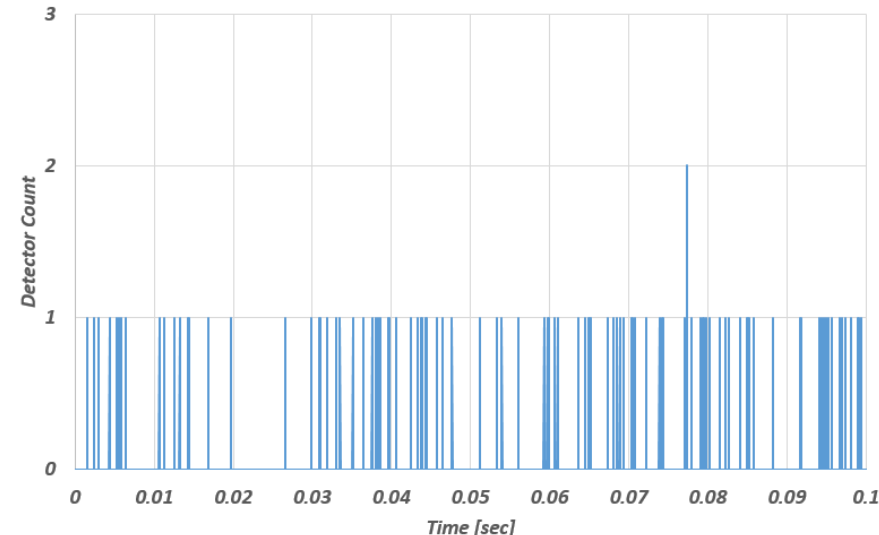
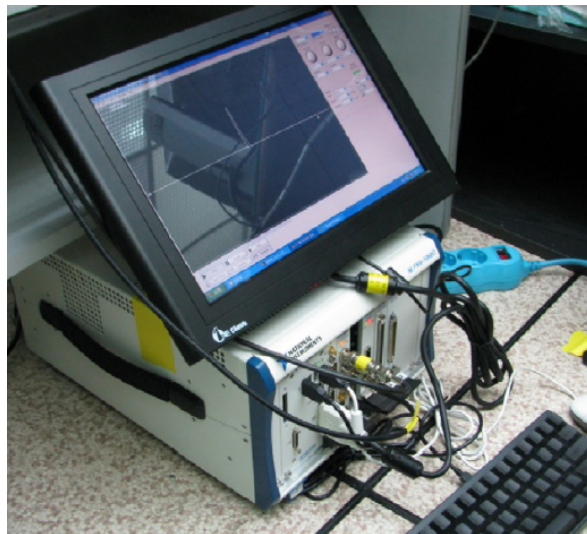


Fig. 3 2nd order Feynman fitting for measured σ_2

❖ *Subcriticality Measurement System (SMS)*

- In this study, a time-series data of neutron counts within a **fine unit gate time of 10 μsec** is acquired using the SMS which was developed by Korea Electric Power Research Institute (KEPRI) for measuring the ex-core detector signal from commercial PWR to get the condition of large subcriticality.
- Since the neutron generation time (Λ) is estimated about 50~60 μsec , **the shorter gate time can acquire more detailed information for estimating α value.**



❖ *Data Processing with a Fully Random Sampling*

- In general, the Feynman- α method requires sufficient number of measurement data for the reliable accuracy of curve fitting.
- A method called “Bunching-technique (time-swap)” increases the number of data by using shifted data even for long gate times.
- However, those method have some disadvantages that the number of the processing data is too big with a fine gate time, which drastically increases computing time.

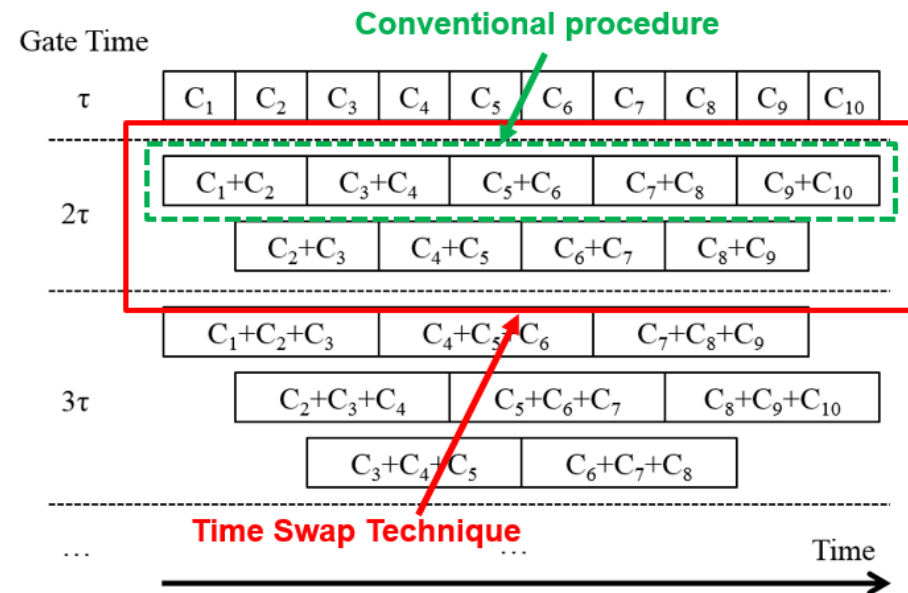
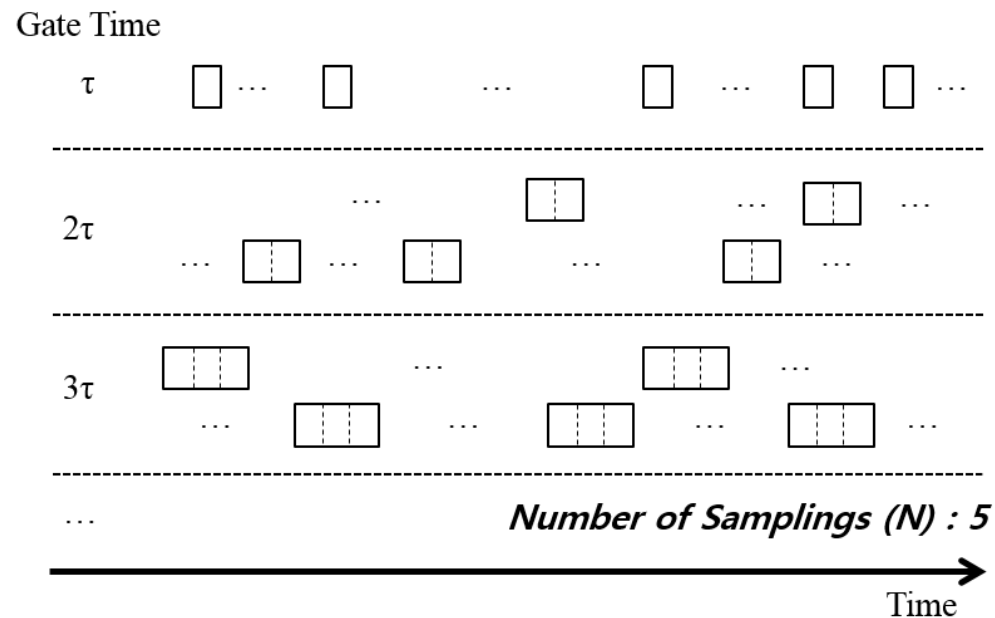


Fig. 5 Schematic view of the bunching-technique

❖ *Data Processing with a Fully Random Sampling*

- In this work, a simple efficient **fully random sampling technique** is suggested to overcome these drawbacks.
- In this method, for a given gate time, a given number of starting time points are randomly sampled over the whole data range and then the consecutive count data within the gate time for each sampled starting time bin : C_k .



❖ Data Processing with a Fully Random Sampling

- The only inputs to be specified are **the length of gate times (or number of gate times)** and the **number of the random samplings** for each time bin.
- As the number of sampling data increases, the measured Y or σ_2 approaches a single value and dispersion decreases.

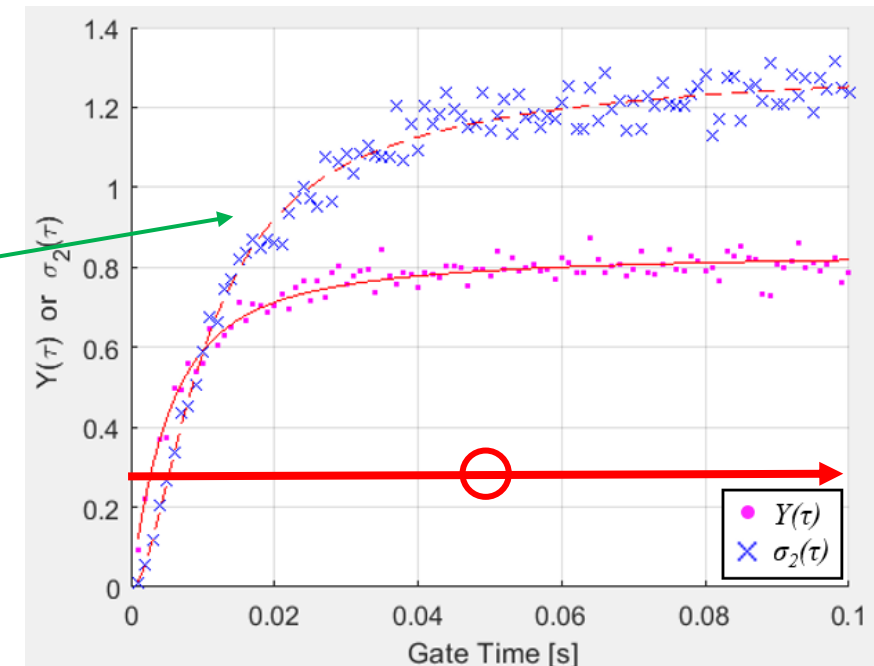
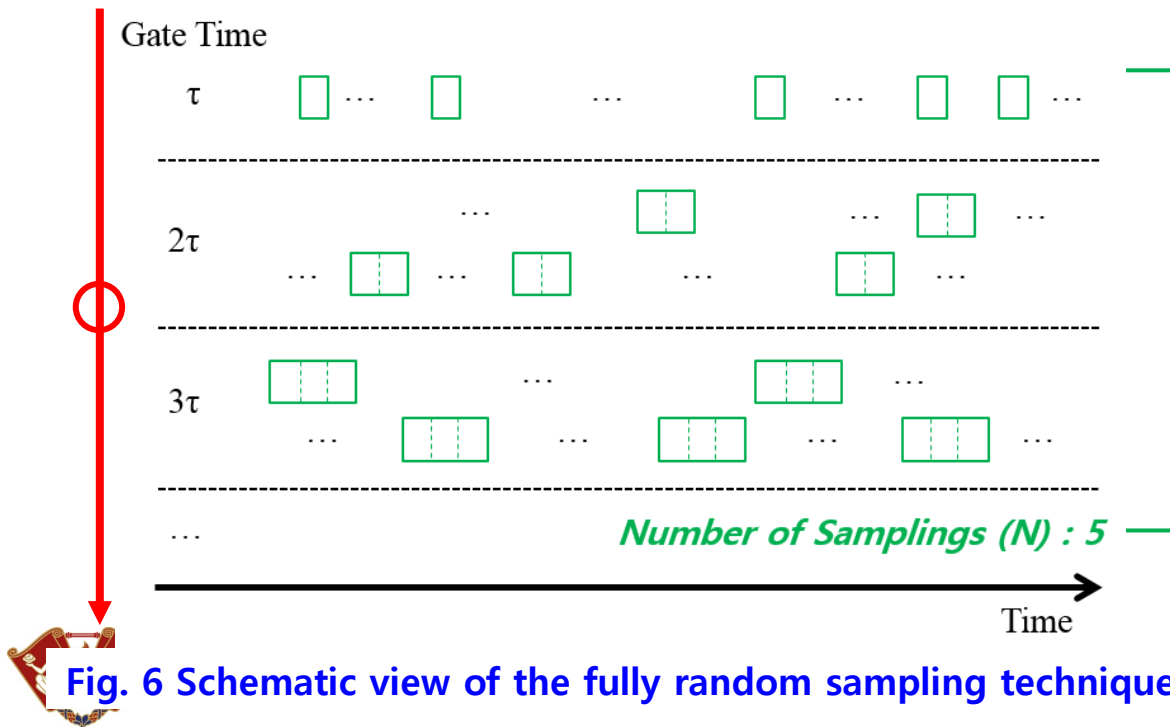


Fig. 7 Feynman fitting using random sampling
(# of gate time times : 100, # of samplings : 10 000)

AGN-201K AND SUBCRITICAL STATES

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❖ AGN-201K

- AGN-201K is a zero-power research and training reactor built by Aerojet General Nucleonics (AGN).
- It is solid moderated reactor using polyethylene and licensed maximum power is 10 Watt.
- The fuel is a **homogeneous mixture of UO₂ and polyethylene**.
- The fuel is comprised of 10 disks with 12.8 cm radius and 25 cm active core height.
- Uranium enrichment of the fuel is about 19.5 w/o.
- The active core is surrounded by 25 cm thick graphite reflector followed by a 10 cm thick lead gamma shield.

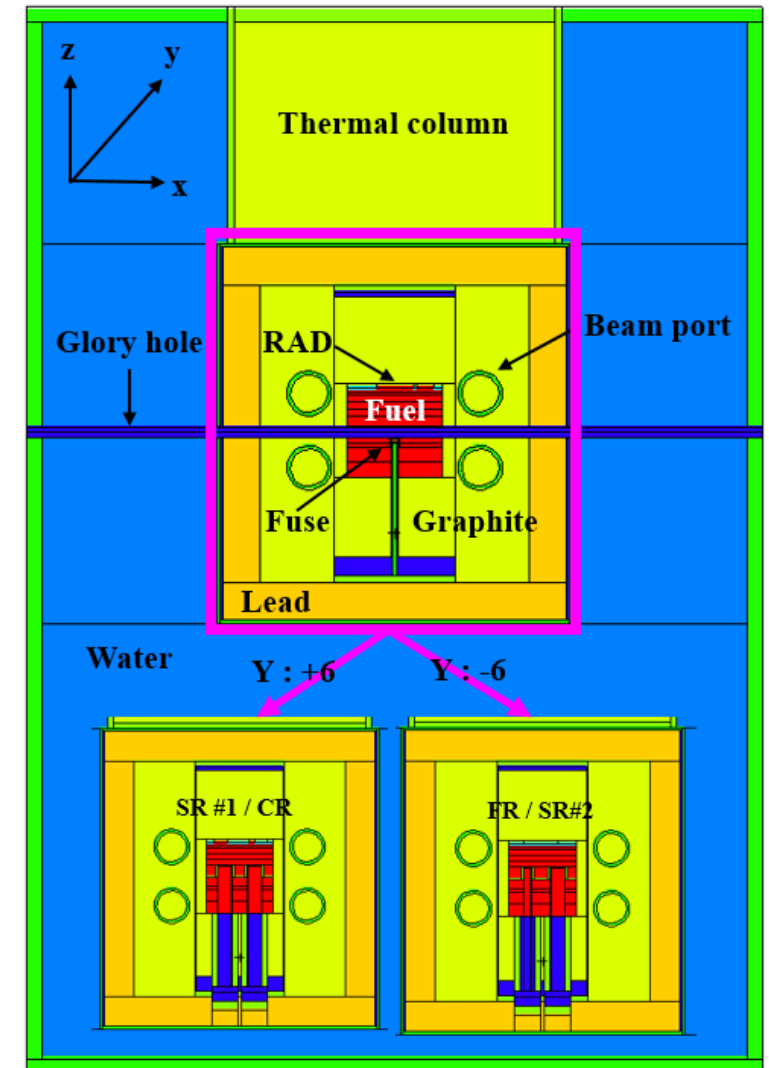


Fig. 8 Axial configuration of the AGN-201K



❖ AGN-201K

- For fast neutron shielding, the outside of the core tank is filled with water of ~47.5 cm thickness.
- The control rod consists of 2 Safety Rods (SR), 1 Coarse Rod (CR), and 1 Fine Rod (FR) that have the same composition as the fuel material.
- During operation reactor power is controlled by CR and FR.
- In particular, an external Ra-Be source located in the left upper beam port supplies neutrons with an intensity of 10 mCi.
- A He-3 ex-core detector connected with SMS is located in right-lower beam port

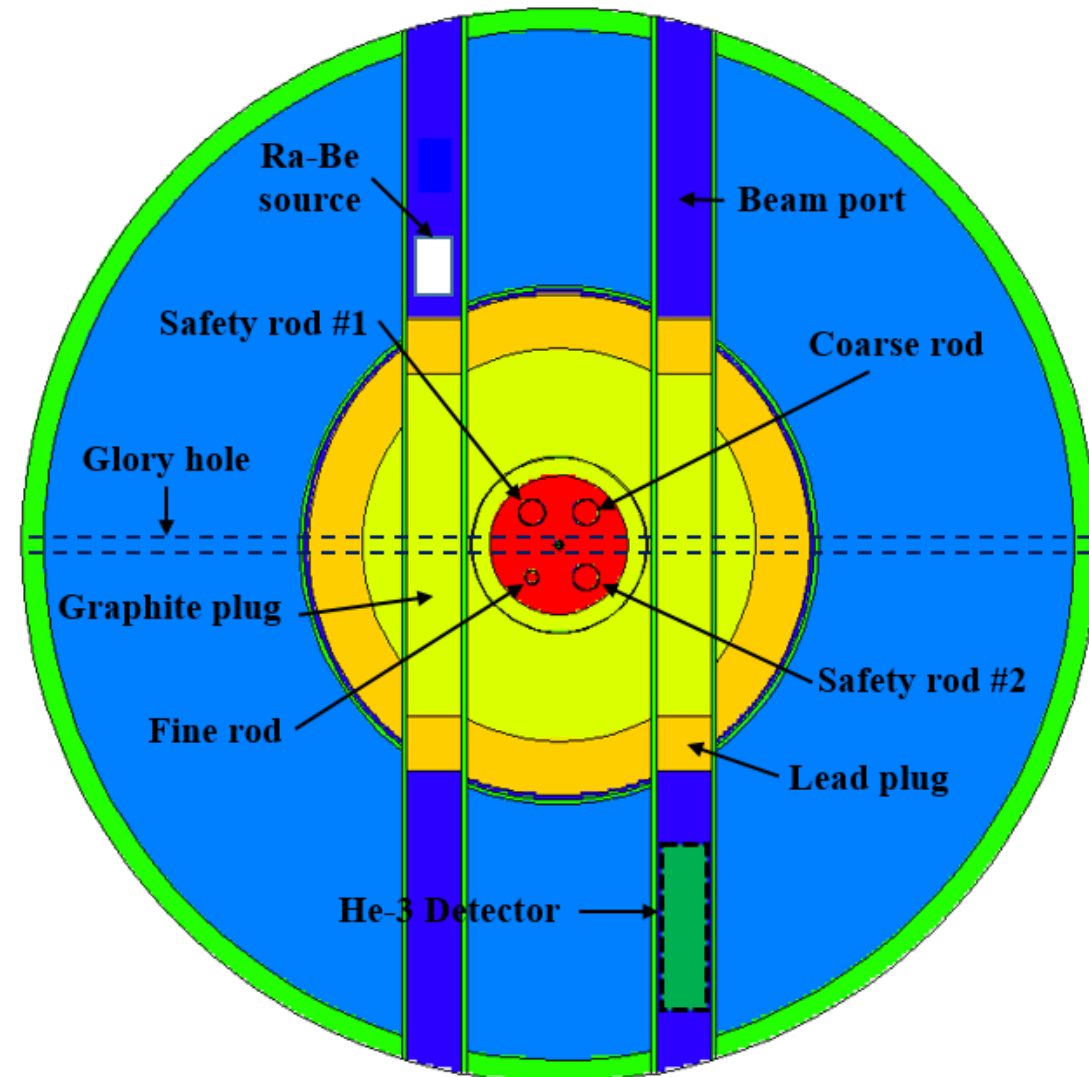


Fig. 10 Radial configuration of the AGN-201K



❖ Selected Subcritical States

- Before the subcriticality measurement, five Sub-Critical condition (SCR) are determined.
- By using the MCNP6 eigenvalue calculations, we obtain the **reference k_{eff}** values and **kinetic parameters**.
- The MCNP6 eigenvalue calculations are performed with ENDF/B-VII.1 cross sections, and with 100 inactive and 5 000 active cycles of 100 000 histories to minimize the statistical error of k_{eff} and kinetic parameters.

TABLE I. Reference multiplication factors and kinetic parameters estimated with MCNP6

Condition	k_{eff}	σ (pcm)	β_{eff}	Λ (μsec)	Inserted rod position (cm)			
					SR#1	SR#2	CR	FR
SCR1	0.98764	3	0.00755	55.89873	23.07	23.44	0	12.56
SCR2	0.99668	3	0.00761	54.55938	23.07	23.44	17.25	12.56
SCR3	0.99737	3	0.00746	54.27638	23.07	23.44	18.25	12.56
SCR4	0.99811	3	0.00757	53.91283	23.07	23.44	19.25	12.56
SCR5	0.99885	3	0.00763	54.00546	23.07	23.44	20.25	12.56

RESULTS AND DISCUSSION

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❖ Experiment Condition

- A series of neutron count was obtained for 5 subcritical conditions by using SMS with a unit gate time of 10 μ sec during 4 minutes to 5 minutes.
- The number of time bins considered was **25 million counts (i.e., 25 000 000 τ , $\tau = 10 \mu$ sec).**
- For curve fitting, the length of gate time was considered up to **0.1 sec (i.e., 10 000 τ).**

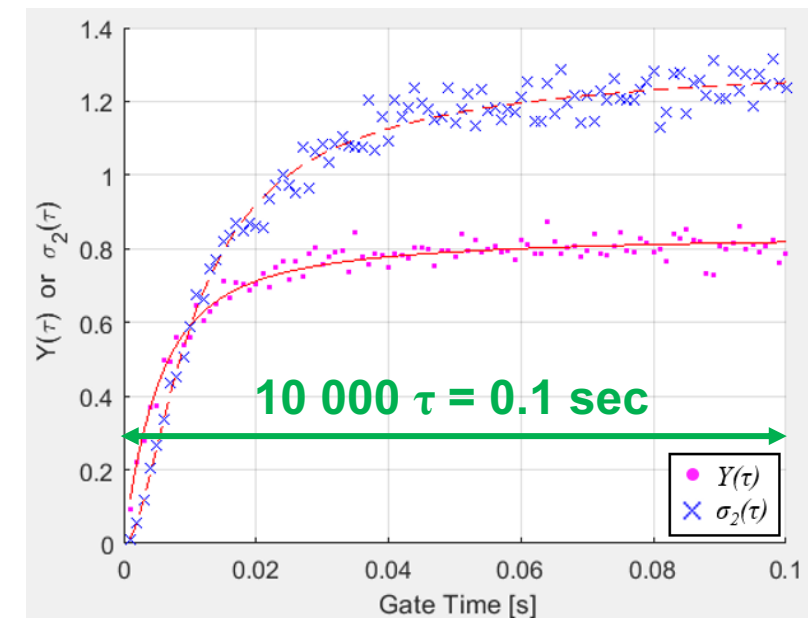
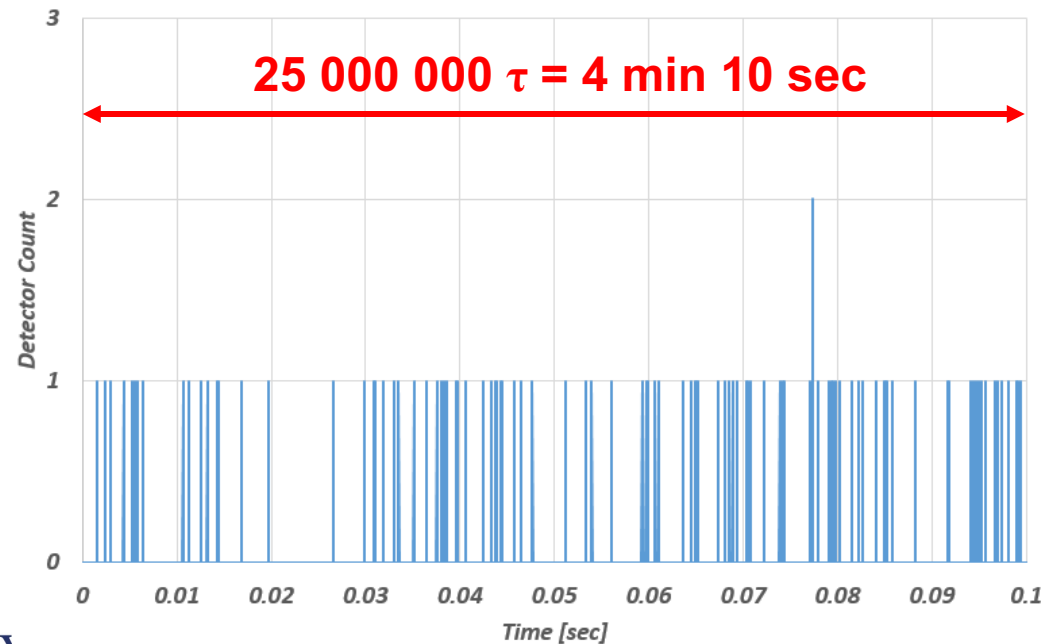


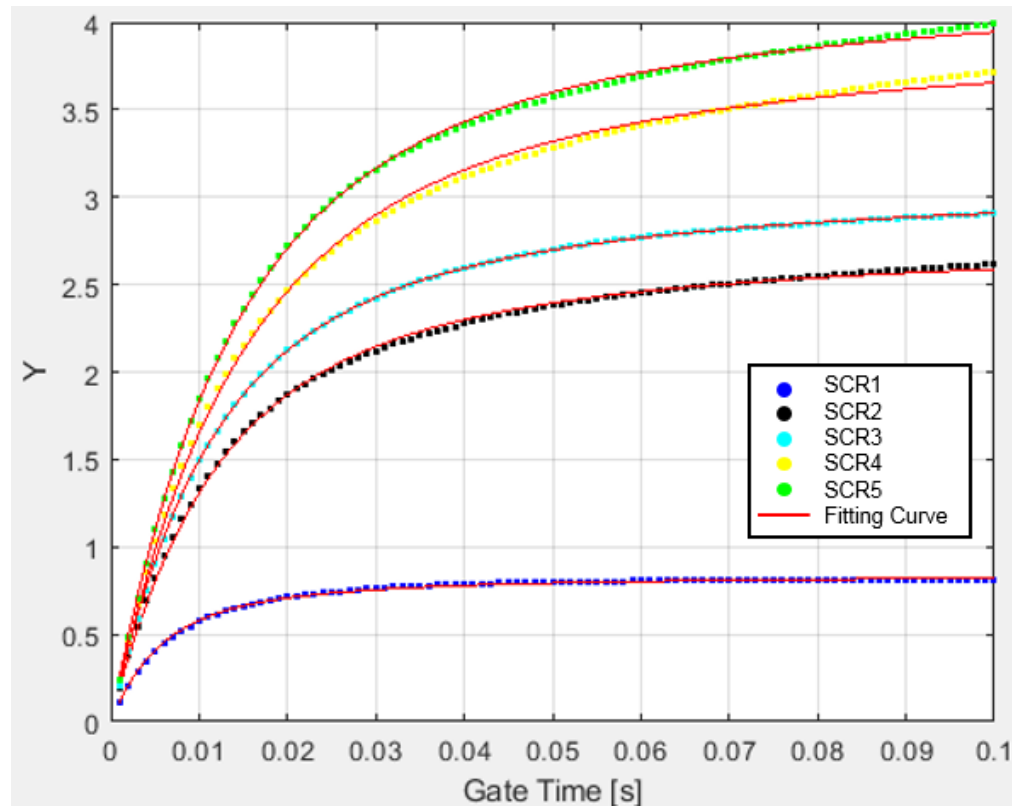
Fig. 11 Experiment condition for measuring subcriticality

RESULTS AND DISCUSSION

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❖ Feynman- α method

- Fig. 12 shows the reference k_{eff} and α value (α -PKE) and difference between reference value and estimated value.



Technique			Time-swap	
# of gate times			100	
# of samplings			whole data	
Condition	k_{eff}	α -PKE	k-est	α -est
SCR1	0.98764	358.95	0.99103 a -338.98	296.99 b 61.96
SCR2	0.99668	200.53	0.99976 -307.80	143.92 56.61
SCR3	0.99737	186.03	0.99940 -202.96	148.51 37.52
SCR4	0.99811	175.53	1.00120 -309.26	118.13 57.40
SCR5	0.99885	162.6	1.00092 -207.39	124.19 38.41
Average CPU time (sec)			108479	

^a $[(k_{\text{eff}}) - (k\text{-est})]$ (pcm Δk), ^b $[(\alpha\text{-PKE}) - (\alpha\text{-est})]$ (1/s)



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Fig. 12 Feynman fitting for five subcritical states using bunching-technique



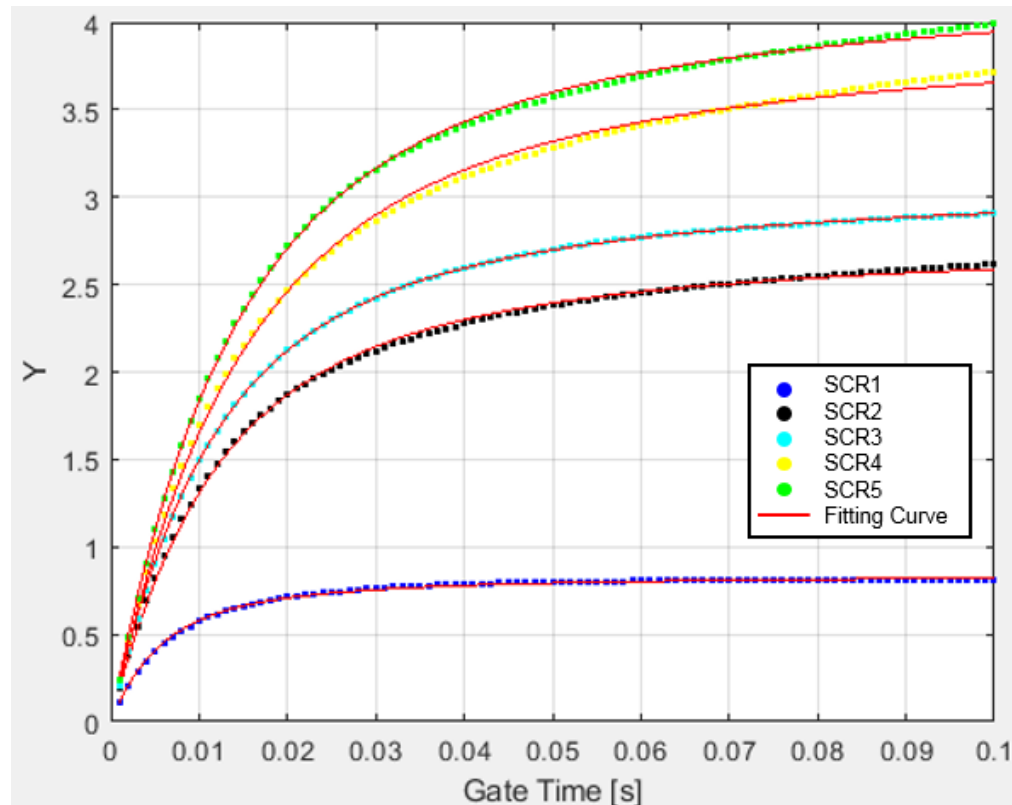
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RESULTS AND DISCUSSION

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❖ Feynman- α method

- Feynman- α method with time swap gives accurate k_{eff} results less than 340 pcm, but long computing times.



Technique			Time-swap	
# of gate times			100	
# of samplings			whole data	
Condition	k_{eff}	α -PKE	k-est	α -est
SCR1	0.98764	358.95	0.99103 a -338.98	296.99 b 61.96
SCR2	0.99668	200.53	0.99976 -307.80	143.92 56.61
SCR3	0.99737	186.03	0.99940 -202.96	148.51 37.52
SCR4	0.99811	175.53	1.00120 -309.26	118.13 57.40
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Average CPU time (sec)			108479	

^a $[(k_{\text{eff}}) - (k\text{-est})]$ (pcm Δk), ^b $[(\alpha\text{-PKE}) - (\alpha\text{-est})]$ (1/s)



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Fig. 12 Feynman fitting for five subcritical states using bunching-technique



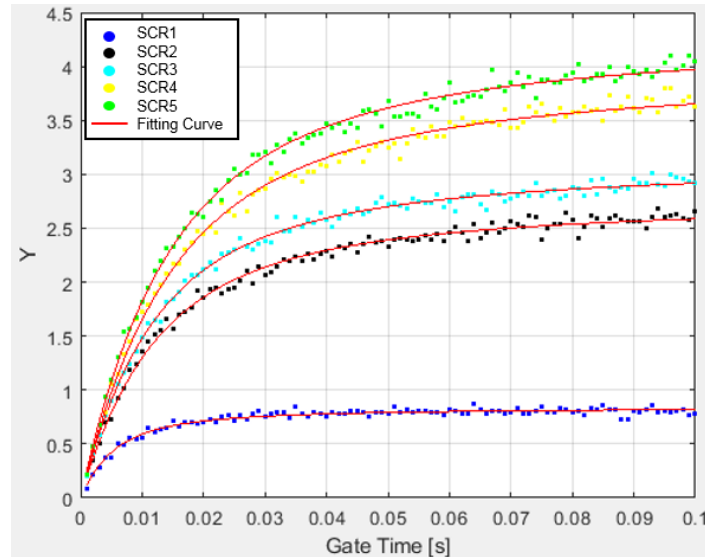
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RESULTS AND DISCUSSION

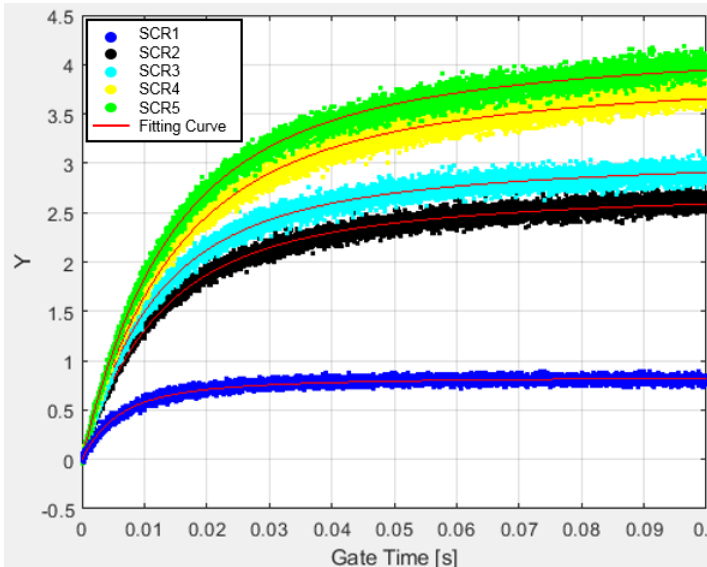
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❖ Feynman- α method

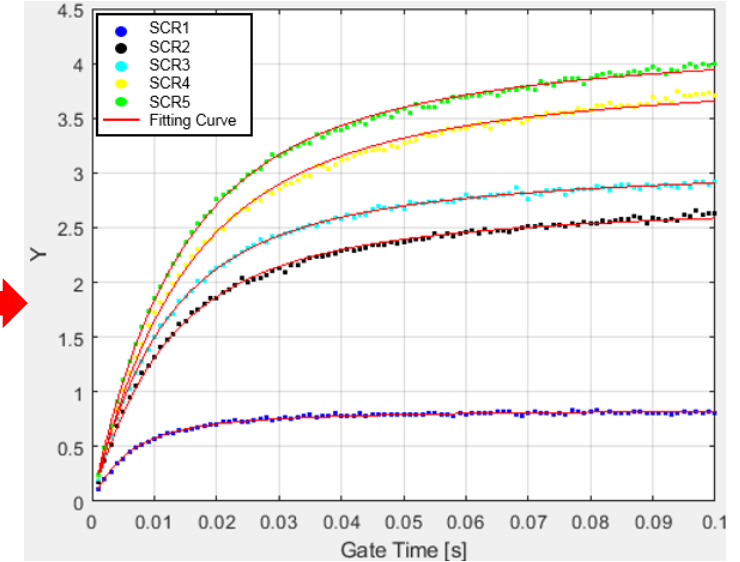
- The largest subcritical state SCR1 shows the lowest slope of the fitting curve.
- As shown in Fig. 13, as the number of sampling increases, the measured Y value converges toward a specific value.



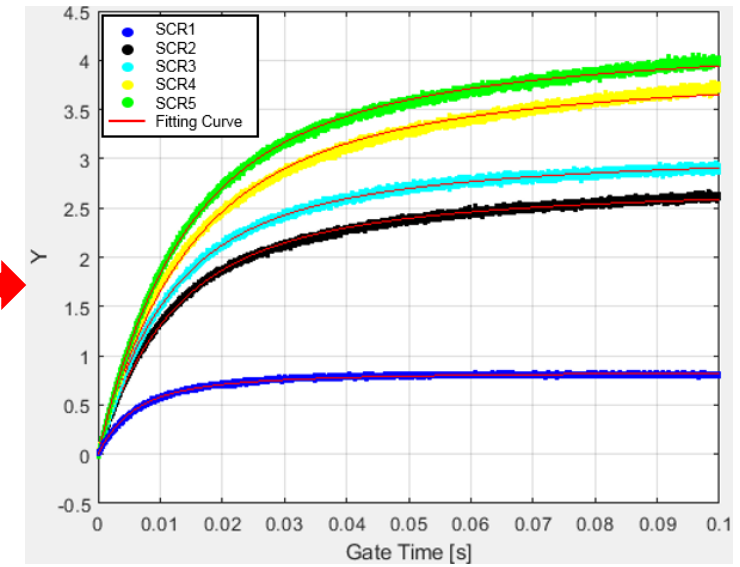
of gate times : 100, # of samplings : 10,000



of gate times : 10,000, # of samplings : 10,000



of gate times : 100, # of samplings : 100,000



of gate times : 10,000, # of samplings : 100,000

Fig. 13 Feynman fitting for five subcritical states using fully random sampling technique

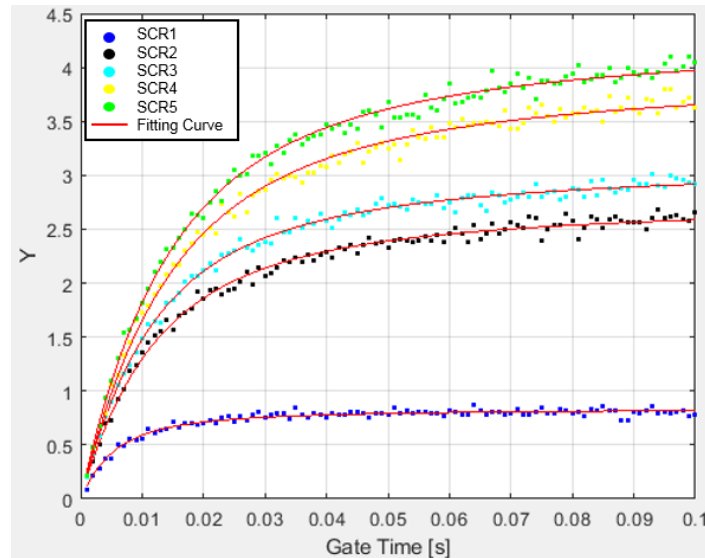


RESULTS AND DISCUSSION

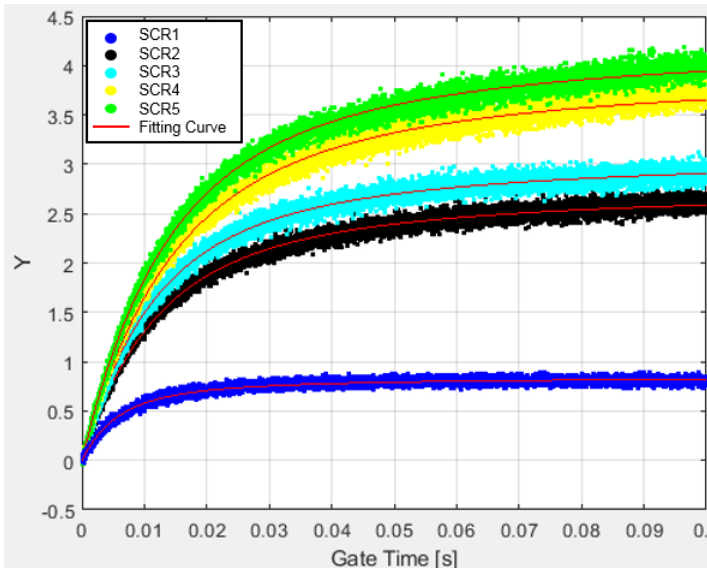
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❖ Feynman- α method

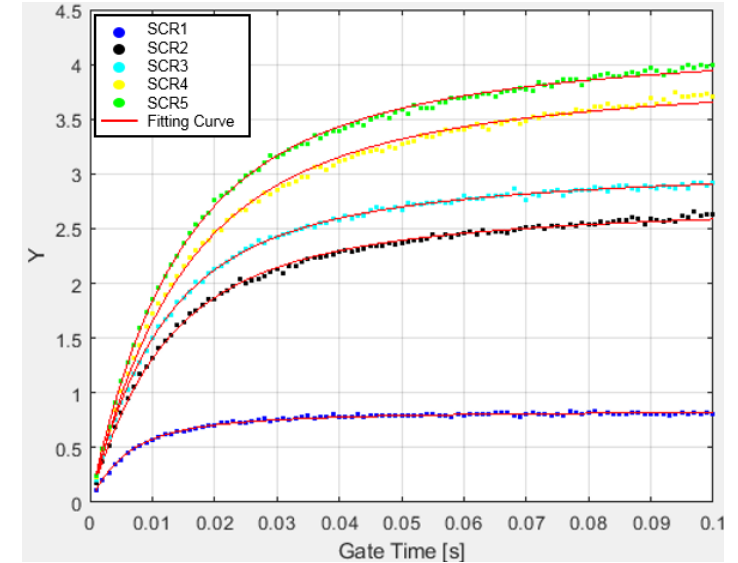
- As the number of gate times increases, the measured α value can be estimated more elaborately.
- Therefore, the last case shows almost the same results as bunching technique even it takes shorter times.



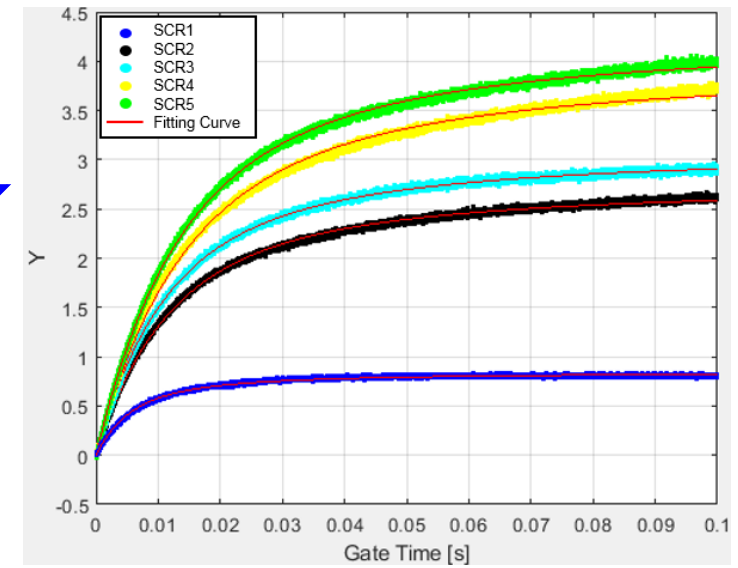
of gate times : 100, # of samplings : 10,000



of gate times : 10,000, # of samplings : 10,000



of gate times : 100, # of samplings : 100,000



of gate times : 10,000, # of samplings : 100,000

Fig. 13 Feynman fitting for five subcritical states using fully random sampling technique



RESULTS AND DISCUSSION

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❖ Feynman- α method with time-swap and fully random sampling methods

- It is noted that fully random sampling method even with a much smaller number of gate times gives comparable accuracies and its computing times are much shorter than those of time-swap method.

TABLE II. Results of Feynman- α method using the time-swap and fully random sampling techniques

Technique			Time-swap		Fully random sampling							
# of gate times			100		100		100		100,000		100,000	
# of samplings			whole data		10,000		100,000		10,000		100,000	
Condition	k _{eff}	α-PKE	k-est	α-est	k-est	α-est	k-est	α-est	k-est	α-est	k-est	α-est
SCR1	0.98764	358.95	0.99103	296.99	0.98985	318.45	0.99115	294.83	0.99102	297.19	0.99103	297.00
			^a -338.98	^b 61.96	-221.31	40.50	-350.84	64.12	-337.86	61.76	-338.95	61.95
SCR2	0.99668	200.53	0.99976	143.92	0.99985	142.27	0.99973	144.35	0.99976	143.80	0.99975	143.98
			-307.80	56.61	-316.80	58.26	-305.46	56.18	-308.45	56.73	-307.45	56.55
SCR3	0.99737	186.03	0.99940	148.51	0.99961	144.61	0.99939	148.69	0.99940	148.51	0.99939	148.64
			-202.96	37.52	-224.12	41.42	-201.99	37.34	-202.95	37.52	-202.29	37.39
SCR4	0.99811	175.53	1.00120	118.13	1.00119	118.36	1.00123	117.62	1.00121	118.01	1.00120	118.23
			-309.26	57.40	-308.01	57.17	-312.03	57.91	-309.91	57.52	-308.74	57.30
SCR5	0.99885	162.6	1.00092	124.19	1.00109	121.10	1.00091	124.36	1.00092	124.25	1.00092	124.19
			-207.39	38.41	-224.10	41.50	-206.49	38.24	-207.06	38.35	-207.41	38.41
Average CPU time (sec)			108479		703		1142		5737		45353	

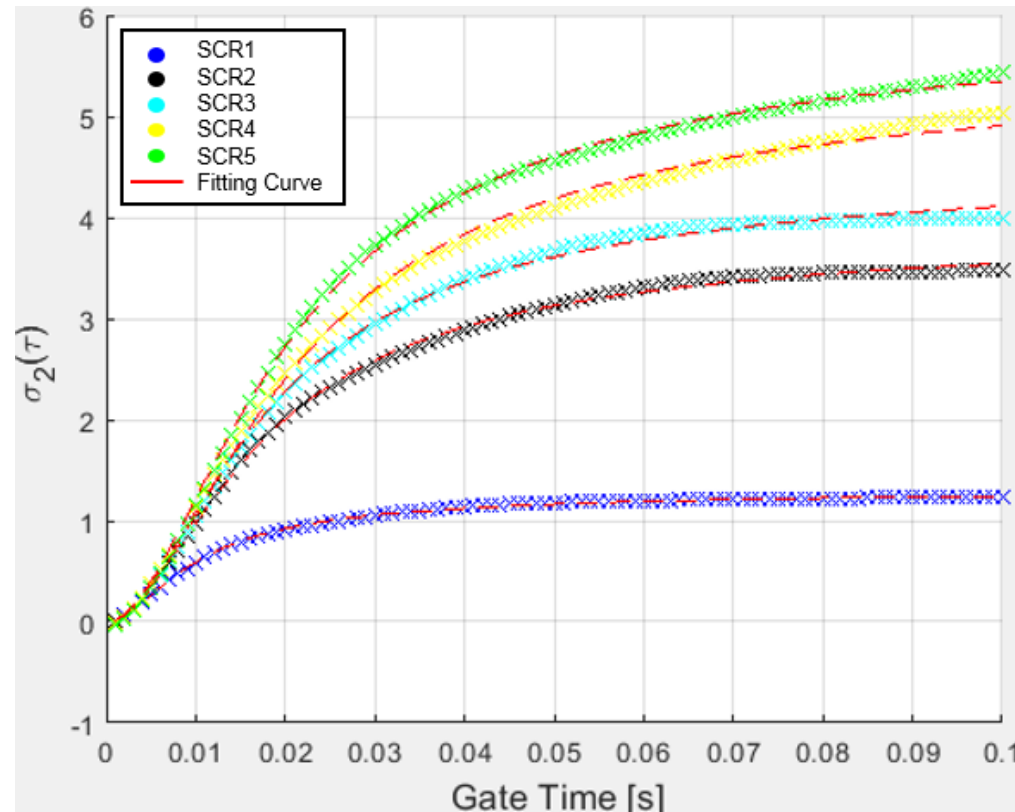
^a $[(k_{\text{eff}}) - (k\text{-est})]$ (pcm Δk), ^b $[(\alpha\text{-PKE}) - (\alpha\text{-est})]$ (1/s)

RESULTS AND DISCUSSION

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❖ 2nd order Feynman- α differential filtering method

- 2nd Feynman- α method with time swap gives accurate k_{eff} results less than 480 pcm, but long computing times.



Technique			Time-swap	
# of gate times			100	
# of samplings			whole data	
Condition	k_{eff}	α -PKE	k-est	α -est
SCR1	0.98764	358.95	0.99247 a -483.34	270.73 b 88.22
SCR2	0.99668	200.53	0.99897 -229.46	158.29 42.24
SCR3	0.99737	186.03	0.99914 -176.95	153.31 32.72
SCR4	0.99811	175.53	1.00059 -247.84	129.50 46.03
SCR5	0.99885	162.6	1.00029 -143.56	136.00 26.60
Average CPU time (sec)			292271	

^a $[(k_{\text{eff}}) - (k\text{-est})]$ (pcm Δk), ^b $[(\alpha\text{-PKE}) - (\alpha\text{-est})]$ (1/s)



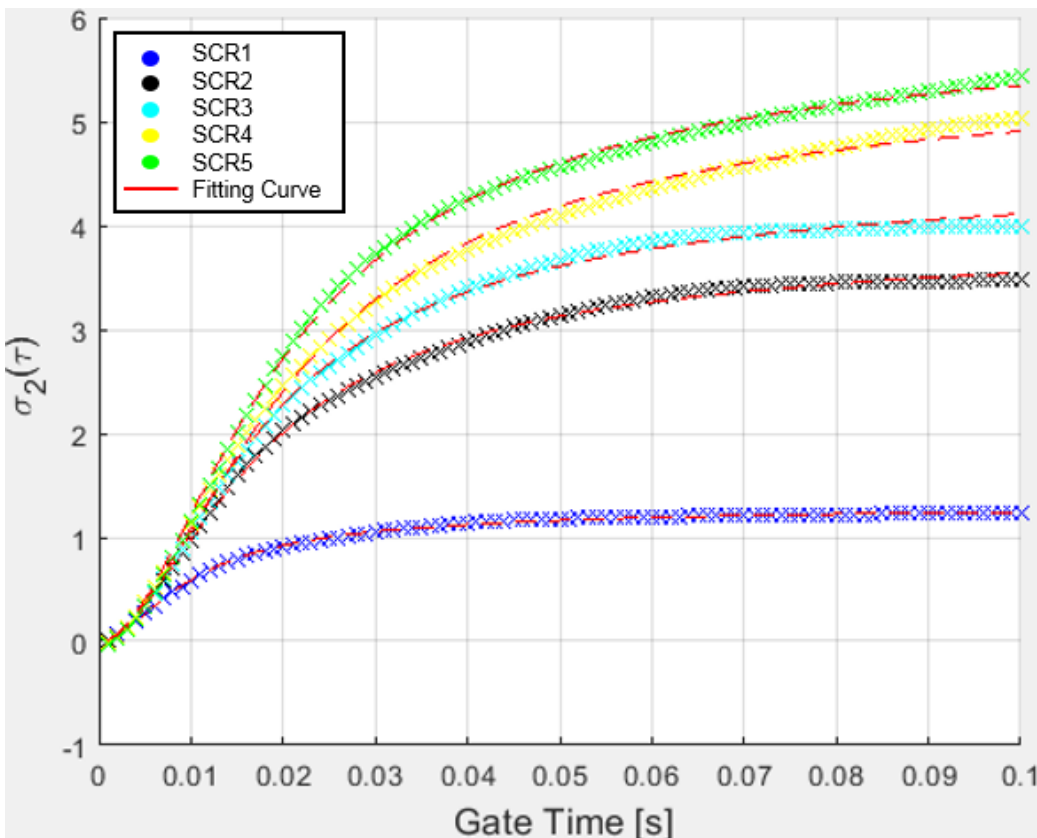
Fig. 14 2nd order Feynman- α fitting for five subcritical states using bunching-technique

RESULTS AND DISCUSSION

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❖ 2nd order Feynman- α differential filtering method

- 2nd Feynman- α method shows more accurate measurement near critical states than conventional Feynman- α method.



Method			2 nd order F- α		Feynman- α	
Technique			Time-swap		Time-swap	
# of gate times			100		100	
# of samplings			whole data		whole data	
Condition	k_{eff}	α -PKE	k-est	α -est	k-est	α -est
SCR1	0.98764	358.95	0.99247 ^a -483.34	270.73 ^b 88.22	0.99103 -338.98	296.99 61.96
SCR2	0.99668	200.53	0.99897 -229.46	158.29 42.24	0.99976 -307.80	143.92 56.61
SCR3	0.99737	186.03	0.99914 -176.95	153.31 32.72	0.99940 -202.96	148.51 37.52
SCR4	0.99811	175.53	1.00059 -247.84	129.50 46.03	1.00120 -309.26	118.13 57.40
SCR5	0.99885	162.6	1.00029 -143.56	136.00 26.60	1.00092 -207.39	124.19 38.41
Average CPU time (sec)			292271		108479	

^a $[(k_{\text{eff}}) - (k\text{-est})]$ (pcm Δk), ^b $[(\alpha\text{-PKE}) - (\alpha\text{-est})]$ (1/s)



Fig. 14 2nd order Feynman- α fitting for five subcritical states using bunching-technique

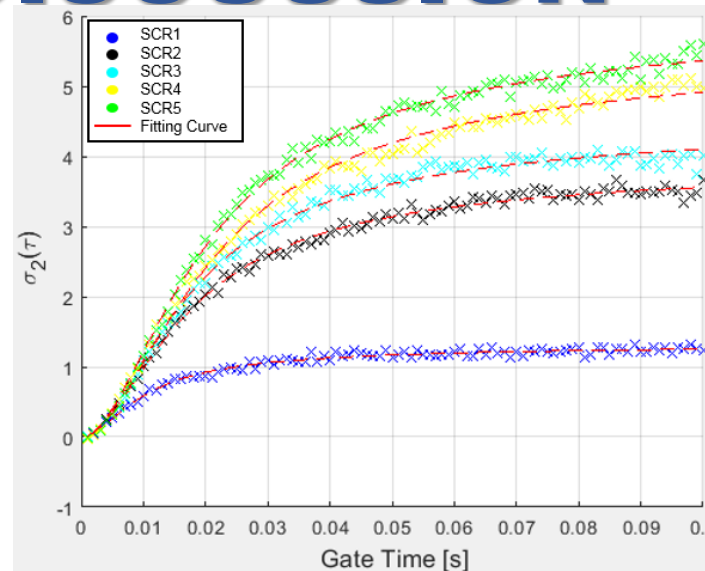
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RESULTS AND DISCUSSION

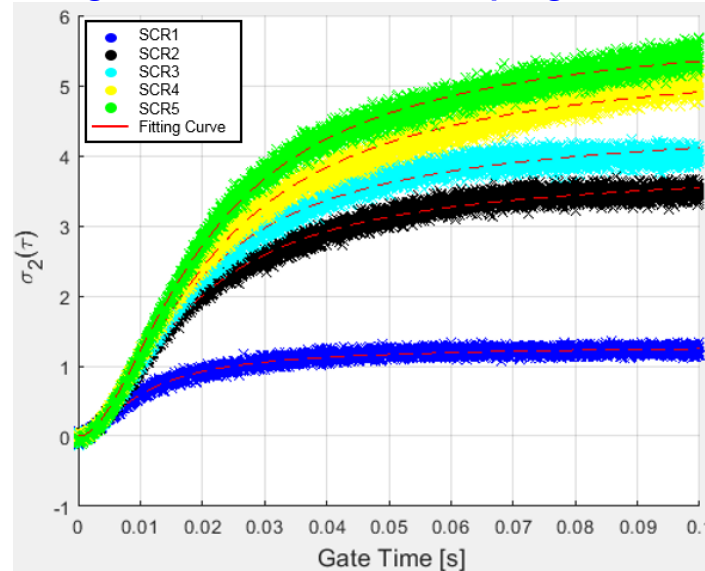
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❖ 2nd Feynman- α method

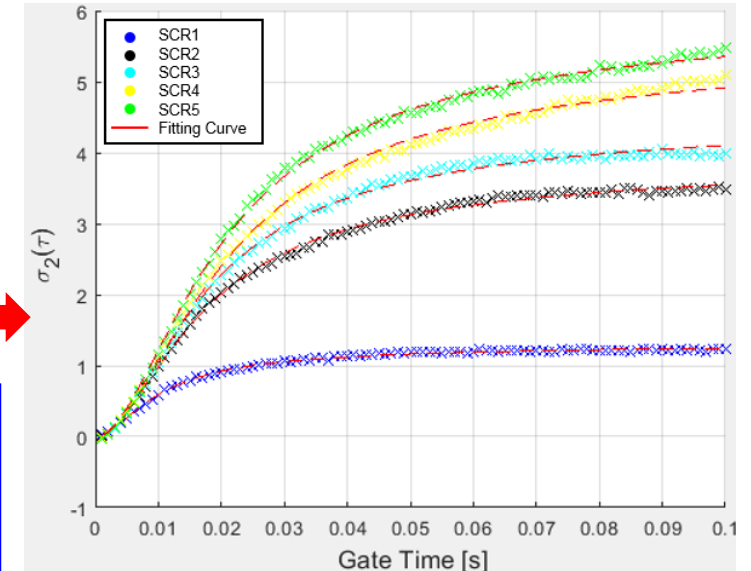
- The shape of the graph looks similar compared with the Feynman- α method, but slight difference at the front.
- As shown in Fig. 15, the number of sampling increases, the measured σ_2 value converges toward a specific value.
- The α value can be estimated more elaborately if we increase the number of gate times.



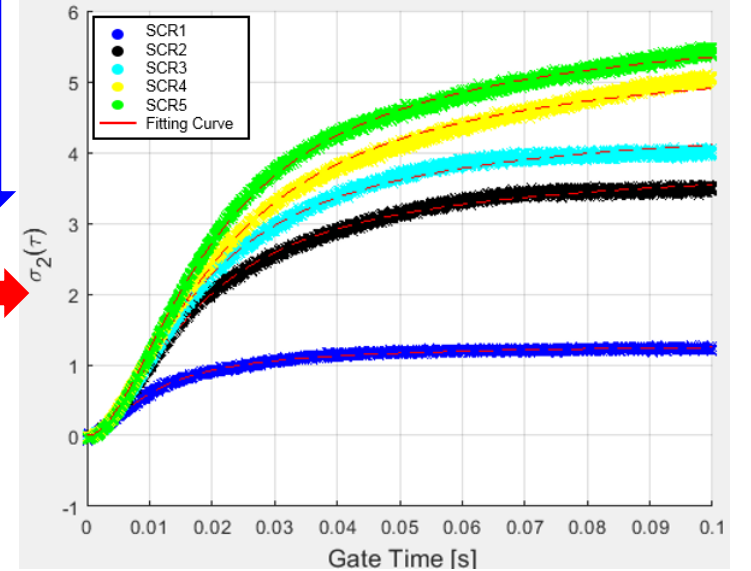
of gate times : 100, # of samplings : 10,000



of gate times : 10,000, # of samplings : 10,000



of gate times : 100, # of samplings : 100,000



of gate times : 10,000, # of samplings : 100,000

Fig. 15 2nd order Feynman- α fitting for five subcritical states using fully random sampling technique



RESULTS AND DISCUSSION

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❖ 2nd order Feynman- α differential filtering method

- It is noted that fully random sampling method even with a much smaller number of gate times gives comparable accuracies and its computing times are much shorter than those of time-swap method.

TABLE III. k_{eff} estimated with 2nd order Feynman- α differential filtering method using the time-swap and fully random sampling techniques

Technique			Time-swap		Fully random sampling							
# of gate times			100		100		100		100,000		100,000	
# of samplings			whole data		10,000		100,000		10,000		100,000	
Condition	k _{eff}	α-PKE	k-est	α-est	k-est	α-est	k-est	α-est	k-est	α-est	k-est	α-est
SCR1	0.98764	358.95	0.99247	270.73	0.99271	266.40	0.99249	270.36	0.99247	270.77	0.99249	270.45
			^a -483.34	^b 88.22	-507.22	92.55	-485.39	88.59	-483.16	88.18	-484.89	88.50
SCR2	0.99668	200.53	0.99897	158.29	0.99897	158.33	0.99896	158.50	0.99897	158.32	0.99898	158.22
			-229.46	42.24	-229.26	42.20	-228.36	42.03	-229.35	42.21	-229.85	42.31
SCR3	0.99737	186.03	0.99914	153.31	0.99906	154.78	0.99913	153.45	0.99915	153.14	0.99914	153.27
			-176.95	32.72	-169.00	31.25	-176.22	32.58	-177.88	32.89	-177.16	32.76
SCR4	0.99811	175.53	1.00059	129.50	1.00057	129.81	1.00060	129.23	1.00060	129.36	1.00058	129.63
			-247.84	46.03	-246.19	45.72	-249.35	46.30	-248.62	46.17	-247.16	45.90
SCR5	0.99885	162.6	1.00029	136.00	1.00031	135.54	1.00030	135.67	1.00029	135.95	1.00028	136.03
			-143.56	26.60	-146.01	27.06	-145.30	26.93	-143.81	26.65	-143.38	26.57
Average CPU time (sec)			292271		719		1401		7224		61796	

^a $[(k_{\text{eff}}) - (k\text{-est})]$ (pcm Δk), ^b $[(\alpha\text{-PKE}) - (\alpha\text{-est})]$ (1/s)

RESULTS AND DISCUSSION

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❖ 2nd order Feynman- α differential filtering method

- It is noted that the 2nd order Feynman- α method shows a more accurate value than the conventional Feynman- α method near critical states (SCR2~SCR5).

TABLE III. k_{eff} estimated with 2nd order Feynman- α differential filtering method using the time-swap and fully random sampling techniques

Method			2 nd order Feynman- α					Conventional Feynman- α				
Technique			Time-swap	Fully random sampling				Time-swap	Fully random sampling			
# of gate times			100	100	100	100,000	100,000	100	100	100	100,000	100,000
# of samplings			whole data	10,000	100,000	10,000	100,000	whole data	10,000	100,000	10,000	100,000
Condition	k _{eff}	α -PKE	k-est	k-est	k-est	k-est	k-est	k-est	k-est	k-est	k-est	k-est
SCR1	0.98764	358.95	0.99247	0.99271	0.99249	0.99247	0.99249	0.99103	0.98985	0.99115	0.99102	0.99103
			^a -483.34	-507.22	-485.39	-483.16	-484.89	^a -338.98	-221.31	-350.84	-337.86	-338.95
SCR2	0.99668	200.53	0.99897	0.99897	0.99896	0.99897	0.99898	0.99976	0.99985	0.99973	0.99976	0.99975
			-229.46	-229.26	-228.36	-229.35	-229.85	-307.80	-316.80	-305.46	-308.45	-307.45
SCR3	0.99737	186.03	0.99914	0.99906	0.99913	0.99915	0.99914	0.99940	0.99961	0.99939	0.99940	0.99939
			-176.95	-169.00	-176.22	-177.88	-177.16	-202.96	-224.12	-201.99	-202.95	-202.29
SCR4	0.99811	175.53	1.00059	1.00057	1.00060	1.00060	1.00058	1.00120	1.00119	1.00123	1.00121	1.00120
			-247.84	-246.19	-249.35	-248.62	-247.16	-309.26	-308.01	-312.03	-309.91	-308.74
SCR5	0.99885	162.6	1.00029	1.00031	1.00030	1.00029	1.00028	1.00092	1.00109	1.00091	1.00092	1.00092
			-143.56	-146.01	-145.30	-143.81	-143.38	-207.39	-224.10	-206.49	-207.06	-207.41
Average CPU time (sec)			292271	719	1401	7224	61796	108479	703	1142	5737	45353

- ❖ In this work, subcriticality experiment is performed with the Feynman- α and 2nd Feynman- α differential method at AGN-201K.
- ❖ A fully random sampling technique is devised to overcome the drawbacks that bunching-technique with fine unit gate time drastically increases computing time,.
- ❖ For measuring subcriticality, eigenvalue calculations are performed with MCNP6 to obtain reference k_{eff} and kinetic parameters.
- ❖ In conclusion, it was shown that the new fully random sampling technique suggested in this work can provide accurate subcriticality estimations with computationally efficient way for AGN-201K and this method coupled with the second-order differential method gives slightly better estimation for the near-critical cases.

**THANK YOU FOR
YOUR ATTENTION !!**



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